Reaching 1 billion rows / second

Hans-Jürgen Schönig www.postgresql-support.de

▲口▶ ▲□▶ ▲目▶ ▲目▶ 三日 ● のへの

Reaching a milestone

・ロト・日本・日本・日本・日本・日本



Traditionally:

- ► We could only use 1 CPU core per query
- Scaling was possible by running more than one query at a time
- Usually hard to do

・ロト・西ト・田・・田・ ひゃぐ



- ▶ PL/Proxy is a stored procedure language to scale out to shards.
- Worked nicely for OLTP workloads
- Somewhat usable for analytics
 - A LOT of manual work

$\mathsf{PL}/\mathsf{Proxy:}$ The future



- Still ok for OLTP
- Certainly not the way to scale out in the future
- Too much manual work
- Not transparent
- Not cool enough

On the app level



Doing scaling on the app level

- A lot of manual work
- Not cool enough
- Needs a lot of development
- Why use a database if work is still manual?
- Solving things on the app level is certainly not an option

The goal

◆□▶ ◆□▶ ◆ 臣▶ ◆ 臣▶ ○ 臣 ○ の Q @



- Import massive amounts of data
- Run typical aggregates
- Process 1 billion rows in less than a second
- Scale out to as many nodes as needed

Coming up with a data structure



▲□▶ ▲圖▶ ▲≣▶ ▲≣▶ = 差 = のへで

We tried to keep that simple:

node=# `	\d	t_demo					
		Table	"pı	ublic.t_dem	10	н	
Column	Ι	Туре	Ι	Collation	I	Nullable	I
	-+-		-+-		-+		-+
id	Ι	serial	Ι		Ι	not null	Ι
grp	Ι	integer	Ι		Ι		I
data	Ι	real	Ι		Ι		Ι
Indexes	:						
"idx_id" btree (id)							

The query



SELECT grp, count(data) FROM t_demo GROUP BY 1;

Single server performance

▲□▶ ▲□▶ ▲目▶ ▲目▶ 目目 のへで



The main questions are:

- How much can we expect from a single server?
- How well does it scale with many CPUs?
- How far can we get?

・ロト・日本・ヨト・ヨー うくぐ

PostgreSQL parallelism



- Parallel queries have been added in PostgreSQL 9.6
 - It can do a lot
 - It is by far not feature complete yet
- Number of workers will be determined by the PostgreSQL optimizer
 - We do not want that
 - We want ALL cores to be at work



 Usually the number of processes per scan is derived from the size of the table

test=# SHOW min_parallel_relation_size ;
min_parallel_relation_size

8MB (1 row)

One process is added if the tablesize triples

Overruling the planner



- We could never have enough data to make PostgreSQL go for 16 or 32 cores.
- Even if the value is set to a couple of kilobytes.
- > The default mechanism can be overruled:

```
test=# ALTER TABLE t_demo
    SET (parallel_workers = 32);
ALTER TABLE
```

Making full use of cores



- How well does PostgreSQL scale on a single box?
- For the next test we assume that I/O is not an issue
 - ► If I/O does not keep up, CPU does not make a difference
 - Make sure that data can be read fast enough.
- Observation: 1 SSD might not be enough to feed a modern Intel chip

Single node scalability (1)





Single node scalability (2)



- We used a 16 core box here
- As you can see, the query scales up nicely
- Beyond 16 cores hyperthreading kicks in
 - ▶ We managed to gain around 18%

Single node scalability (3)



- On a single Google VM we could reach close to 40 million rows / second
- For many workloads this is already more than enough
- Rows / sec will of course depend on type of query

Moving on to many nodes

▲□▶ ▲□▶ ▲目▶ ▲目▶ 目目 のへで



- We want to shard data to as many nodes as needed
- ▶ For the demo: Place 100 million rows on each node
 - \blacktriangleright We do so to eliminate the I/O bottleneck
 - In case I/O happens we can always compensate using more servers
- Use parallel queries on each shard



explain SELECT grp, COUNT(data) FROM t_demo GROUP BY 1; Finalize HashAggregate

Group Key: t_demo.grp

- -> Append
 - -> Foreign Scan (partial aggregate)
 - -> Foreign Scan (partial aggregate)
 - -> Partial HashAggregate
 - Group Key: t_demo.grp
 - -> Seq Scan on t_demo





- Throughput doubles as long as partial results are small
- Planner pushes down stuff nicely
- Linear increases are necessary to scale to 1 billion rows

Preconditions to make it work (1)



- postgres_fdw uses cursors on the remote side
 - cursor_tuple_fraction has to be set to 1 to improve the planning process
 - set fetch_size to a large value
- That is the easy part



- We have to make sure that all remote nodes work at the same time
- This requires "parallel append and async fetching"
 - All queries are sent to the many nodes in parallel
 - Data can be fetched in parallel



- PostgreSQL could not be changed without substantial work being done recently
 - Traditionally joins had to be done BEFORE aggregation
 - This is a showstopper for distributed aggregation because all the data has to be fetched from the remote host before aggregation
- Kyotaro Horiguchi fixed, which made our work possible
 - This was a HARD task !



Easy tasks:

- Aggregates have to be implemented to handle partial results coming from shards
- Code is simple and available as extension

・ロト・日本・ヨト・ヨー うくぐ



- Dissect aggregation
- Send partial queries to shards in parallel
- Perform parallel execution on shards
- Add up data on main node

Final results



Hardware used



- We used 32 boxes (16 cores) on Google
- Data was in memory
- Adding more servers is EASY

Future ideas



- JIT compilation will speed up execution
- More parallelism for more executor nodes
- General speedups (tuple deforming, etc.)
- In the future FEWER cores will be needed to achieve similar results

Contact us



Cybertec Schönig & Schönig GmbH Hans-Jürgen Schönig Gröhrmühlgasse 26 A-2700 Wiener Neustadt

www.postgresql-support.de

Follow us on Twitter: @PostgresSupport